

DADE GROUP

12.4.7 - 12.4.15

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THE POINT: INTRODUCE A CONSTRUCTION THAT
CAN BE USED TO RECOVER A LOT OF THE
OPERATIONS WE CAN DO ON $\text{Perm}_K(P)$
LIKE Res , Iso ,

12.4.7 GENERALIZED TENSOR INDUCTION

LET P AND Q BE FINITE p -GROUPS AND
 U BE A FINITE (Q, P) -BISET.

IF X IS A FINITE P -SET, DENOTE BY

$$T_U(X) := \text{Hom}_P(U^{\circ P}, X) = \left\{ \varphi \mid \varphi: U \rightarrow X \text{ s.t. } \varphi(ug^{-1}) = g\varphi(u), \forall g \in P, \forall u \in U \right\}$$

\parallel

$$\varphi(g \cdot_P u)$$

NOTE: $T_U(X) = \text{Hom}_P(U^{\text{op}}, X)$ IS A FINITE Q-SET
WITH THE FOLLOWING ACTION!

$$\text{GIVEN } \varphi \in T_U(X), h \in Q \quad (h \cdot \varphi)(u) = \varphi(h^{-1}u)$$

$\forall u \in U.$

NEED TO CHECK $(h \cdot \varphi) \in \text{Hom}_P(U^{\text{op}}, X)$?

$$\text{NOTE } (h \cdot \varphi)(u \cdot \bar{g}) \stackrel{\text{by } Q}{=} \varphi(h^{-1}u \bar{g}) \stackrel{\text{by } P}{=} g \cdot (\varphi(h^{-1}u))$$

$$\stackrel{\text{by } Q}{=} g \cdot ((h \cdot \varphi)(u)) \text{ so } (h \cdot \varphi) \in T_U(X), \forall g \in P.$$

EXAMPLE:

LET $Q \subseteq P$, AND) LET $V = P$ AS (Q, P) -BASED, X P -SET

NOTE $T_U(X) = \text{Hom}_P(P, X) \cong \text{Res}_Q^P X$ } ON OBJECTS
 $\alpha \mapsto \alpha(1) \in X$

$$(h \cdot \alpha)(1) \underset{\substack{\downarrow \\ \text{by } Q}}{=} \alpha(h^{-1}) \underset{\substack{\downarrow \\ \text{by } P}}{=} h \cdot (\alpha(1))$$

NOTE T_U GIVE US A MAP:

$$T_U : \text{Obj}(\underline{\text{Perm}}_K(P)) \longrightarrow \text{Obj}(\underline{\text{Perm}}_K(Q))$$

$$X \longmapsto T_U(X)$$

CAN WE HAVE A COMPATIBLE MAP ON MORPHISM?

LET $m \in \text{Hom}_{\underline{\text{Perm}}_K(P)}(X, Y)$ SO RECALL

$$m : Y \times X \rightarrow K \quad \text{s.t.} \quad m(g_Y, g_X) = m(Y, X) \quad \text{for all } g \in P$$

$$\forall (Y, X) \in Y \times X$$

WANT $T_U(m) \in \text{Hom}_{\text{PER}_{m \times m}(\mathbb{Q})}(T_U(X), T_U(Y))$

DEFINE:

$$T_U(m) : T_U(Y) \times T_U(X) \rightarrow K$$
$$(\psi, \varphi) \mapsto T_U(m)(\psi, \varphi) = \prod_{u \in [U/\rho]} m(\overset{x}{\psi}(u), \overset{y}{\varphi}(u))$$

WHERE $[U/\rho]$ IS A SET OF REPRESENTATIVES OF U/ρ .

WELL DEFINED? $m(\varphi(u g^{-1}), \varphi(u g^{-1})) = m(g \varphi(u), g \varphi(u))$

WITH $g \in P$.

$$m(\varphi(u), \varphi(u))$$

AS m, φ, φ ARE P -EQUIVARIANT.

12.4.8. Lemma : This definition yields a morphism $T_U(m)$ from $T_U(X)$ to $T_U(Y)$ in the category $\text{Perm}_k(Q)$.

PROOF: NEED TO SHOW $T_U(m)$ IS Q -INVARIANT, $h \in Q$

$$T_U(m)(h^{-1}\varphi, h^{-1}\varphi) \stackrel{\textcircled{1}}{=} \prod_{u \in [U/P]} m(h^{-1}\varphi(u), h^{-1}\varphi(u)) \text{ by DEF } T_U(m)$$

$$\text{by DEF OF } Q \text{ ACTION} \stackrel{\textcircled{2}}{=} \prod_{u \in [U/P]} m(\varphi(h \cdot u), \varphi(hu))$$

THE IMAGE OF all $h \cdot u$ on $[U/P]$

IS ANOTHER SET OF REPRESENTATIVES.

$$\begin{aligned} &= \pi_{u' \in [U/\rho]} m(\varphi(u'), \varphi(u')) = T_U(m)(\varphi, \varphi) \end{aligned}$$

SO THE MATRIX $T_U(m)$ IS \mathbb{Q} -INVARIANT.

SO $T_U(m)$ IS A MORPHISM IN $\underline{\text{Perm}}_K(Q)$.

THUS ONE CAN ASK IS $T_U(m)$ A FUNCTOR?

12.4.9. Lemma : [19, Lemma 2.2] Let k be a field of characteristic p . Let P and Q be finite p -groups, and let U be a finite (Q, P) -biset. Then T_U is a functor from $\text{Perm}_k(P)$ to $\text{Perm}_k(Q)$.

PROOF: LEFT TO SHOW: ① $T_U(\text{Id}_X) = \text{Id}_{T_U(X)}$

$$\textcircled{2} T_U(m, m) = T_U(m) \circ T_U(m)$$

① LET X BE FINITE P -SET, $\text{Id}_X = \delta_X(y, y') = \begin{cases} 1, & y = y' \\ 0, & y \neq y' \end{cases}$

LET $\varphi, \varphi' \in T_U(X)$

$$T_U(\delta_X)(\varphi, \varphi') = \prod_{u \in U/P} \delta_X(\varphi(u), \varphi'(u)) = \begin{cases} 0, & \varphi \neq \varphi' \\ 1, & \varphi = \varphi' \end{cases}$$

So ① ✓. AS ψ, φ ARE P-EQUIVARIANT

② ... LEFT TO THE READER...

NEED TO SHOW $m \mapsto T_U(m)$ IS MULTIPLICATIVE

$$T_U(m \circ n)(\psi, \theta) = T_U(m) \circ T_U(n)(\psi, \theta)$$

WHERE $m: Z \rightarrow X$ AND $n: X \rightarrow Y$

$$\theta \in T_U(Z), \quad \psi \in T_U(Y)$$

NOTE: UNDER CERTAIN CONDITIONS THERE ARE MAPS FROM $F: \text{Perm}_k(P) \rightarrow \text{Perm}_k(Q)$. SO IT IS NATURAL TO ASK IS THERE A \cup A FINITE (Q,P) -BISET MAKES THE DIAGRAM COMMUTES OR THE OTHER

(WAY AROUND)

$$\begin{array}{ccc}
 \underline{\text{Perm}}_k(P) & \xrightarrow{e_P} & \text{Perm}_k(P) \\
 \downarrow T_U & & \downarrow F \\
 \underline{\text{Perm}}_k(Q) & \xrightarrow{e_Q} & \text{Perm}_k(Q)
 \end{array}$$

12.4.10 EXAMPLE: ELEMENTARY BISSETS:

EXAMPLE ①: F-RESTRICTION

LET $Q \leq P$, A^m) LET $U = P$ as (Q, P) -biset, X P -SET

NOTE $T_U(X) = \text{Hom}_P(P, X) \cong \text{Res}_Q^P X$ } ON OBJECTS
 $\alpha \mapsto \alpha(1) \in X$

LET $m: X \rightarrow Y$ IN $\text{PER}_{m_X}(P)$ THEN } ON MORPHISM

$$T_U(m)(\varphi, \varphi) = \prod_{u \in P/P} m(\varphi(u), \varphi(u)) = m(\varphi(1), \varphi(1))$$

$$\text{So } T_U(m) = m$$

So T_U IS JUST THE RESTRICTION FUNCTOR.

$$\begin{array}{ccc} \underline{\text{Perm}}_K(P) & \xrightarrow{e_P} & \text{Perm}_K(P) \\ T_U \downarrow & & \downarrow \text{Res}_P^P \\ \underline{\text{Perm}}_K(Q) & \xrightarrow{e_Q} & \text{Perm}_K(Q) \end{array}$$

EXAMPLE 2: TENSOR INDUCTION:

LET $Q \subseteq P$, AND) LET $V = P$ AS (P, Q) -BISSET, X Q -SET

$$T_V(X) = \text{Hom}_Q(V^{op}, X) = \text{Hom}_Q(P, X) \cong \prod_{S \in P/Q} X$$

$$\left(\begin{array}{l} \varphi \longmapsto \prod_{S \in P/Q} (\varphi(s)) \\ \varphi \longleftarrow \prod_{S \in P/Q} X_S \end{array} \right)$$

$$\text{by } \varphi(sh) = h^{-1} X_S, \quad h \in Q$$

NOTE: UNDER THE ISOMORPHISM $\text{Hom}_{\mathbb{Q}}(P, X) \cong \prod_{\sigma \in P/Q} X$

ONE CAN GIVE AN ACTION OF $g \in P$ ON $\prod_{\sigma \in P/Q} X$

ONE CAN CHECK: $g \cdot (X_{\sigma})_{\sigma \in S} = (h_{\epsilon^{-1}(\sigma)} X_{\epsilon^{-1}(\sigma)})_{\sigma \in P/Q}$

by $g \cdot \sigma = \epsilon(\sigma) h_{\sigma}$ WHERE ϵ IS A PERMUTATION OF P/Q .
 $h_{\sigma} \in \mathbb{Q}$.

ON ANOTHER NOTE; GIVEN $K \times$ WE CAN DEFINE

$\text{Ten}_Q^P(KX) = \bigotimes_{S \in P/Q} S \otimes KX$ CALLED THE TENSOR INDUCED
MODULE

AS BEFORE, $\mathfrak{g} \left(\bigotimes_{S \in P/Q} S \otimes m_s \right) = \bigotimes_{S \in P/Q} \left(S \otimes h_{\bar{E}(S)}^{-1} m_{\bar{E}(S)} \right)$

WHERE $\mathfrak{g} \in P$,

SO NOTE $K T_V(X) \cong K \pi_X \cong \bigotimes_{S \in P/Q} \text{Ten}_Q^P(KX)$

ONE CAN CHECK!

$$\begin{array}{ccc} \underline{\text{Perm}}_k(Q) & \xrightarrow{e_Q} & \text{Perm}_k(Q) \\ \downarrow T_V & & \downarrow \text{Ten}_Q^P \\ \underline{\text{Perm}}_k(P) & \xrightarrow{e_P} & \text{Perm}_k(P) \end{array}$$

EXAMPLE 3: INFLATION:

LET $N \trianglelefteq P$, $U = P/N$ AS $(P, P/N)$ -~~B~~BISET

LET X BE A P/N -SET

$$T_U(X) = \text{Hom}_{P/N}(P/N, X) \cong \text{Inf}_{P/N}^P X$$

$$\varphi \mapsto \varphi(1)$$

NOTE $T_U(m) = m$ AS $U/P/N = P/N/P/N = 1$.

(SAME ARGUMENT WITH Res)

SO T_U IS JUST THE INFLATION FUNCTOR

$$\begin{array}{ccc}
 \underline{\text{Perm}}_k(P/N) & \xrightarrow{e_{P/N}} & \text{Perm}_k(P/N) \\
 \downarrow T_U & & \downarrow \text{Inf}_{P/N}^P \\
 \underline{\text{Perm}}_k(P) & \xrightarrow{e_P} & \text{Perm}_k(P)
 \end{array}$$

EXAMPLE 4: BRAUER QUOTIENT:

LET $V = P/N$ AS $(P/N, P)$ -BISSET THEN WE HAVE

$$T_V(X) \cong X^{\mathcal{N}} \quad \text{AND} \quad kX^{\mathcal{N}} \cong (kX)[\mathcal{N}] \quad \text{SO}(\dots)$$

||

$\text{Hom}_P(P/N, X)$

$$\begin{array}{ccc}
 \underline{\text{Perm}}_k(P) & \xrightarrow{e_P} & \text{Perm}_k(P) \\
 \downarrow T_V & & \downarrow \text{Br}_N^P \\
 \underline{\text{Perm}}_k(P/N) & \xrightarrow{e_{P/N}} & \text{Perm}_k(P/N)
 \end{array}$$

EXAMPLE 5: ISOMORPHISM:

LET $F: P \xrightarrow{\sim} Q$, X BE A FINITE P -SET

$U = P \rightsquigarrow (Q, \rho)$ -BISSET

$T_U(X) \cong \text{Iso}_P^Q(X)$ SO ONE CAN SHOW:

$$\begin{array}{ccc} \underline{\text{Perm}}_K(P) & \xrightarrow{e_P} & \text{Perm}_K(P) \\ T_U \downarrow & & \downarrow \text{Iso}_P^Q \\ \underline{\text{Perm}}_K(Q) & \xrightarrow{e_Q} & \text{Perm}_K(Q) \end{array}$$

SO T_U IS JUST ISO.

NOTE: RESTRICTION, TENSOR INDUCTION, INFLATION,

BRAUER QUOTIENT, ISOMORPHISM CAN BE RECOVERED BY

MEANS OF THIS T_U . HENCE ^{WE HAVE} THE FOLLOWING DEFINITION

BECAUSE OF TENSOR INDUCTION!

DEFINITION: THE FUNCTOR T_U IS CALLED

THE GENERALIZED TENSOR INDUCTION ASSOCIATED TO U .

SOME PROPERTIES OF T_U !

12.4.12. Proposition : *Let k be a field of characteristic p , let P and Q be finite p -groups, and let U be a finite (Q, P) -biset.*

1. *If X is a one element P -set, then $T_U(X)$ is a one element Q -set.*
2. *If X and Y are finite P -sets, there is an isomorphism*

$$T_U(X \times Y) \cong T_U(X) \times T_U(Y),$$

which is functorial in X and Y : there is a commutative diagram of categories and functors

$$\begin{array}{ccc} \underline{\text{Perm}}_k(P) \times \underline{\text{Perm}}_k(P) & \xrightarrow{\times} & \underline{\text{Perm}}_k(P) \\ T_U \times T_U \downarrow & & \downarrow T_U \\ \underline{\text{Perm}}_k(Q) \times \underline{\text{Perm}}_k(Q) & \xrightarrow{\times} & \underline{\text{Perm}}_k(Q). \end{array}$$

3. *If X and Y are finite P -sets, let $\sigma : X \times Y \rightarrow Y \times X$ denote the switch morphism. Then up to the isomorphism of Assertion 2, the morphism $T_U(\sigma)$ is the switch morphism $T_U(X) \times T_U(Y) \rightarrow T_U(Y) \times T_U(X)$.*
4. *If U' is another finite (Q, P) -biset, then the functors $T_{U \sqcup U'}$ and $T_U \times T_{U'}$ are isomorphic.*

PROOF: ① IF $|X|=1$ THEN THERE IS A UNIQUE MORPHISM OF P-SETS FROM ANY P-SET TO X.

②

2. If X and Y are finite P-sets, there is an isomorphism

$$T_U(X \times Y) \cong T_U(X) \times T_U(Y),$$

THERE IS A CANONICAL ISOMORPHISM OF Q-SETS

$$T_U(X \times Y) = \text{Hom}_P(U^{\text{op}}, X \times Y) \cong \text{Hom}_P(U^{\text{op}}, X) \times \text{Hom}_P(U^{\text{op}}, Y)$$
$$\begin{array}{ccc} & \parallel & \parallel \\ & T_U(X) & \times & T_U(Y) \end{array}$$

ON MORPHISM: GIVEN $m: X \rightarrow X'$ AND $n: Y \rightarrow Y'$

WE DEFINE $m \times n: X \times Y \rightarrow X' \times Y'$ i.e

$$(m \times n)(x', y', x, y) = m(x', x) n(y', y).$$

ALL WE NEED TO SHOW IS:

$$T_U(m \times n): T_U(X) \times T_U(Y) \rightarrow T_U(X') \times T_U(Y')$$

||

$$T_U(m) \times T_U(n): T_U(X) \times T_U(Y) \rightarrow T_U(X') \times T_U(Y')$$

by showing: $T_U(m \times n)(\varphi', \varphi) = T_U(m)(\varphi'_{X'}, \varphi'_X) \cdot T_U(n)(\varphi'_{Y'}, \varphi'_Y)$

$$\varphi': U^{\text{op}} \rightarrow X' \times Y', \quad \varphi: U^{\text{op}} \rightarrow X \times Y$$

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3. If X and Y are finite P -sets, let $\sigma : X \times Y \rightarrow Y \times X$ denote the switch morphism. Then up to the isomorphism of Assertion 2, the morphism $T_U(\sigma)$ is the switch morphism $T_U(X) \times T_U(Y) \rightarrow T_U(Y) \times T_U(X)$.

THE SWITCH MORPHISM $\sigma : X \times Y \rightarrow Y \times X$ IS DEFINED BY

$$\sigma(y, x, x', y') = \delta_{x, x'} \delta_{y, y'}$$

THEN $T_U(\sigma) : (T_U(Y) \times T_U(X)) \times (T_U(X) \times T_U(Y)) \rightarrow K$

$$T_U(\sigma)(\varphi, \psi, \varphi', \psi') = \prod_{u \in U/P} \sigma(\varphi(u), \psi(u), \varphi'(u), \psi'(u))$$

AS $\varphi, \psi, \varphi', \psi'$ ARE P -EQUIVARIANT, FOR $T_U(\sigma)$ TO BE NON-ZERO NEED $\varphi = \varphi'$ AND $\psi = \psi'$.

So $T_U(\sigma)$ IS THE SWITCH MORPHISM.

④

4. If U' is another finite (Q, P) -biset, then the functors $T_{U \sqcup U'}$ and $T_U \times T_{U'}$ are isomorphic.

ON OBJECTS: LET X BE A P -SET

$$T_{U \sqcup U'}(X) = \text{Hom}_P((U \sqcup U')^{op}, X) \cong \underset{\parallel}{\text{Hom}_P(U^{op}, X)} \times \underset{\parallel}{\text{Hom}_P(U'^{op}, X)}$$
$$T_U(X) \quad \times \quad T_{U'}(X)$$

GIVEN BY $(\varphi: (U \sqcup U')^{op} \rightarrow X) \mapsto (\varphi|_U, \varphi|_{U'})$

ON MORPHISM: ^{LET} $m: Y \times X \rightarrow K$ BE A MORPHISM FROM
X TO Y. WANT TO SHOW $T_{U \times U'}(m) \cong T_U(m) \times T_{U'}(m)$

$$T_{U \times U'}(m)(\varphi, \varphi') = \prod_{u \in [U \times U']/\rho} m(\varphi(u), \varphi'(u)) \quad \text{by DEF } T$$

$$\text{by U.} \quad \cong \prod_{u \in [U]/\rho} m(\varphi(u), \varphi'(u)) \cdot \prod_{u' \in [U']/\rho} m(\varphi(u), \varphi'(u'))$$

$$\text{by } T \quad \cong T_U(m)(\varphi|_U, \varphi|_U) \cdot T_{U'}(m)(\varphi|_{U'}, \varphi|_{U'})$$

$$\text{So } T_{U \times U'}(m) \cong T_U(m) \times T_{U'}(m) \quad \square$$

RECAP: SOME PROPERTIES OF T_U !

12.4.12. Proposition : *Let k be a field of characteristic p , let P and Q be finite p -groups, and let U be a finite (Q, P) -biset.*

1. *If X is a one element P -set, then $T_U(X)$ is a one element Q -set.*
2. *If X and Y are finite P -sets, there is an isomorphism*

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which is functorial in X and Y : there is a commutative diagram of categories and functors

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3. *If X and Y are finite P -sets, let $\sigma : X \times Y \rightarrow Y \times X$ denote the switch morphism. Then up to the isomorphism of Assertion 2, the morphism $T_U(\sigma)$ is the switch morphism $T_U(X) \times T_U(Y) \rightarrow T_U(Y) \times T_U(X)$.*
4. *If U' is another finite (Q, P) -biset, then the functors $T_{U \sqcup U'}$ and $T_U \times T_{U'}$ are isomorphic.*

12.4.13 COMPOSITION AND GALOIS TWISTS:

HOW DO WE COMPOSE THOSE OPERATIONS?

LET $P, Q,$ AND R BE FINITE p -GROUPS,

LET U BE A FINITE (Q, P) -BISET AND

V BE A FINITE (R, Q) -BISET.

$V \times_Q U$ BE A FINITE (R, P) -BISET.

$$T_U: \underline{\text{Perm}}_K(P) \rightarrow \underline{\text{Perm}}_K(Q) \quad T_V: \underline{\text{Perm}}_K(Q) \rightarrow \underline{\text{Perm}}_K(R)$$

SO WE CAN FORM: $T_V \circ T_U: \underline{\text{Perm}}_K(P) \rightarrow \underline{\text{Perm}}_K(R)$

AND ALSO $T_{V \times_Q U}: \underline{\text{Perm}}_K(P) \rightarrow \underline{\text{Perm}}_K(R)$

ARE $T_V \circ T_U$ AND $T_{V \times_Q U}$ ISOMORPHIC?

SPOILER ALERT: no

NEED GALOIS TWISTS IN ORDER TO UNDERSTAND
THE CONNECTION BETWEEN $T_V \circ T_U$ AND $T_{V \times_Q U}$.

12.4.14 DEFINITION: LET α BE AN ENDOMORPHISM OF THE FIELD K . LET P BE A FINITE p -GROUP

$$\gamma_{\alpha, P}(\gamma_{\alpha}): \text{Per}_{m, K}(P) \rightarrow \text{Per}_{m, K}(P)$$

$$\text{OBJECTS: } X \mapsto X$$

$$\text{MORPHISMS: } m \mapsto \alpha(m) \quad \left| \begin{array}{l} m: Y \times X \mapsto K \end{array} \right.$$

$$\text{S.T. } \alpha(m)(y, x) = \alpha(m(y, x)).$$

WHY IS γ_a A FUNCTOR?

$$\textcircled{1} \alpha(\text{Id}_X) = \alpha(\delta_X) = \delta_X$$

FOLLOW AS α IS AN

$$\textcircled{2} \alpha(m) \circ \alpha(n) = \alpha(m \circ n) \quad \text{ENDOMORPHISM OF FIELD } K.$$

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FOR $m \in \mathbb{N}$, LET $\gamma(p^m)$ DENOTE THE FUNCTOR γ_a FOR THE ENDOMORPHISM $\alpha: X \rightarrow X^{p^m}$ OF K .

THE FUNCTOR $\gamma(p)$ IS CALLED THE FROBENIUS TWIST FUNCTOR

SOME PROPERTIES OF THE γ_a FUNCTOR:

12.4.15. Lemma :

1. Let a and b be endomorphisms of k . Then for any finite p -groups P

$$\gamma_{a,P} \circ \gamma_{b,P} \cong \gamma_{a \circ b, P} .$$

2. Let P and Q be finite p -groups, and U be a finite (Q, P) -biset. Then

$$T_U \circ \gamma_{a,P} = \gamma_{a,Q} \circ T_U .$$

PROOF: ① ON OBJECTS IT IS JUST THE IDENTITY MAP
ON MORPHISM FOLLOWS BY COMPOSITION OF
ENDOMORPHISM

$$\gamma_{a,P} \circ \gamma_{b,P}(m) = a(\gamma_{b,P}(m)) = a(b(m)) = (a \circ b)(m)$$

②

2. Let P and Q be finite p -groups, and U be a finite (Q, P) -biset. Then

$$T_U \circ \gamma_{a,P} = \gamma_{a,Q} \circ T_U.$$

ON OBJECTS: $T_U \circ \gamma_{a,P}(x) = T_U(x) = \gamma_{a,Q} \circ T_U(x)$

γ IS JUST IDENTITY ON OBJECTS

ON MORPHISM: $m: Y \times X \rightarrow K$

$$(T_U \circ \gamma_{a,P}(m))(\varphi, \psi) = \prod_{u \in U/P} a(m)(\varphi(u), \psi(u)) \quad \text{by DEF. } \gamma$$

$$= \prod_{u \in U/P} a(m(\varphi(u), \psi(u))) \quad \text{by DEF } a$$

$$\cong \alpha \left(\prod_{u \in U/P} m(\varphi(u), \psi(u)) \right)$$

$$= \alpha \left(T_U(m) (\varphi, \psi) \right) \quad \text{by DEF. T}$$

$$= (\gamma_{\alpha, \beta} \circ T_U(m)) (\varphi, \psi) \quad \text{by DEF. } \alpha$$

□

SUMMARY

- DEFINED T_U , FUNCTOR
- SAW SOME RES, ALSO, THE RANGE RECOVER BY T_U
- SAW SOME PROPERTIES OF T_U
- ARE $T_U \circ T_U$ AND $T_U \times_{\mathbb{Q}} U$ ISOMORPHIC?
- NO, NEED NEW FUNCTOR \mathcal{J}_a
- SAW SOME PROPERTIES OF \mathcal{J}_a